

Bugs Framework (BF)

Formalizing Cybersecurity Weaknesses and Vulnerabilities

Office of the National Cyber Director (ONCD)

April 4, 2024

- Introduction
 - CWE, CVE, NVD
 - BF Approach
 - BF Security Concepts
- BF
 - Bugs Models
 - Weakness Taxonomies
 - Vulnerability Models
 - Formal Language
- BF Datasets
 - BFCWE
 - BCVE
- BF Vulnerability Classification Model
- Potential Impacts

Introduction

Current State of the Art

- Weaknesses

[CWE](https://cwe.mitre.org/) – Common Weakness Enumeration

<https://cwe.mitre.org/>

- Vulnerabilities

[CVE](https://cve.mitre.org/) – Common Vulnerabilities and Exposures

<https://cve.mitre.org/>

- Assigning weaknesses to vulnerabilities – CWEs to CVEs

[NVD](https://nvd.nist.gov/) – National Vulnerabilities Database

<https://nvd.nist.gov/>

Repository Challenges

- Imprecise descriptions
- Unclear causality
- Gaps in coverage
- Overlaps in coverage
- Wrong NVD assignments
- No tracking methodology
- No tools

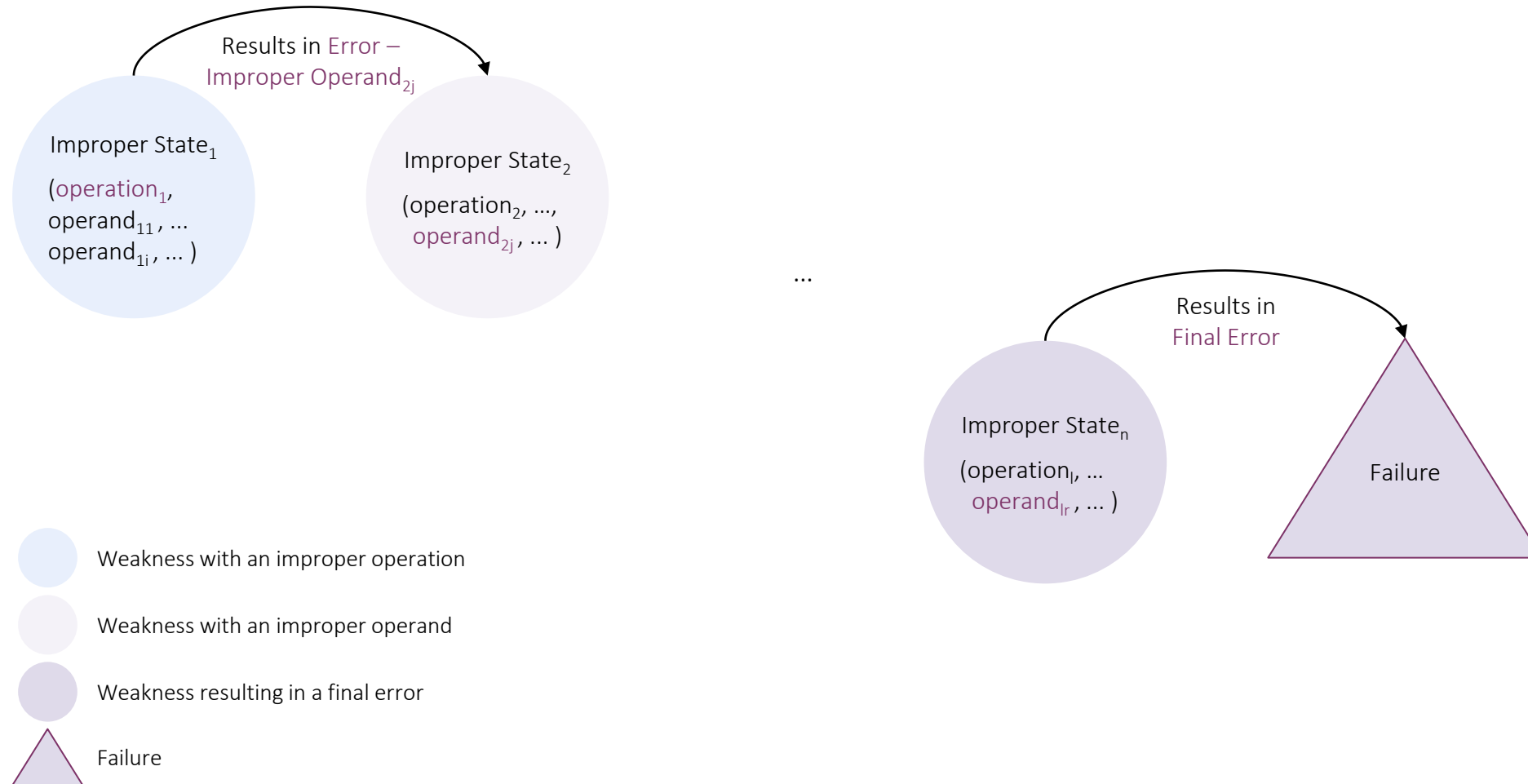
★ Focus is on descriptions

Challenges	Imprecise Descriptions	Unclear Causality	Gaps in Coverage	Overlaps in Coverage	Wrong CVE to CWE mapping	No Tracking Methodology	No Tools
CWE	✓	✓	✓	✓		✓	✓
CVE	✓	✓				✓	✓
NVD	✓	✓			✓	✓	✓

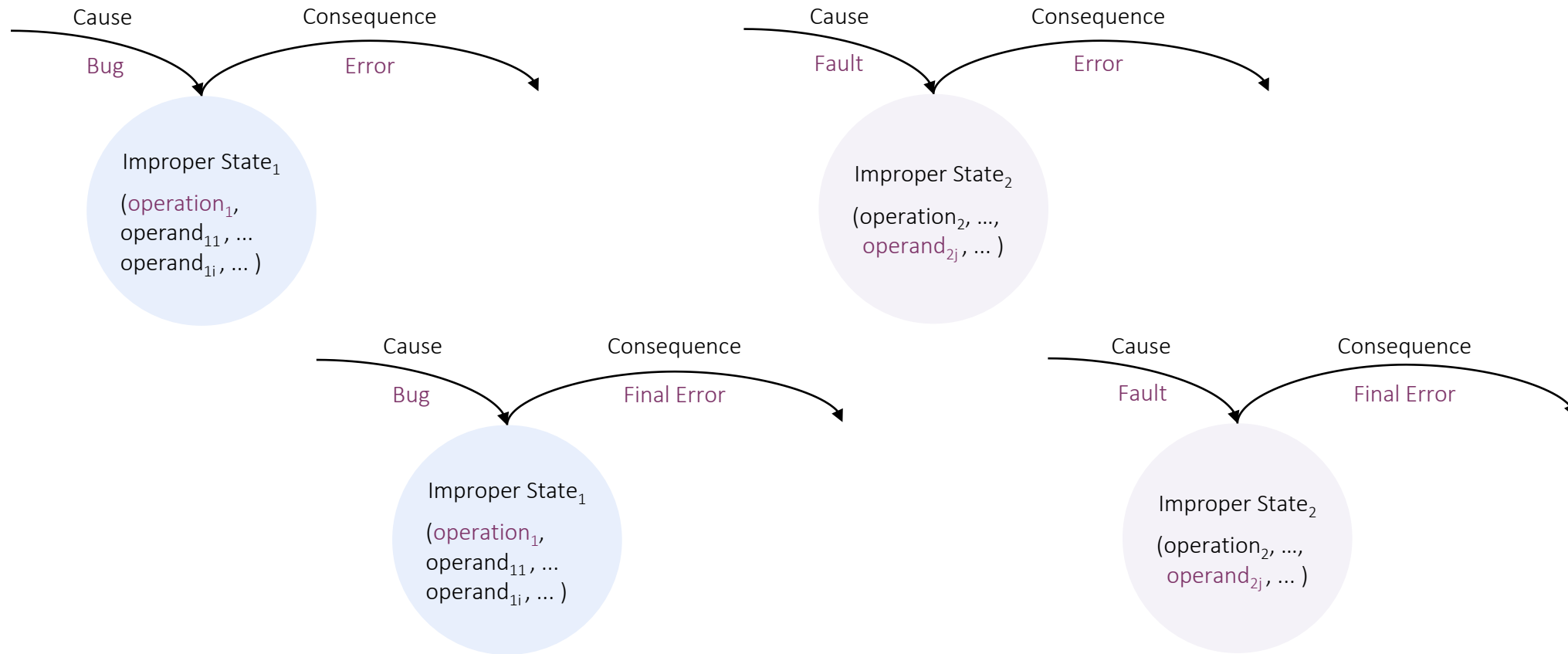
BF is a classification of security bugs and related faults, featuring a formal language for unambiguous specification of weaknesses and underlined by them vulnerabilities.



- Bugs and faults – as weakness causes
- Errors and final errors – as weakness consequences
- BF formal language – based on:
 - Weakness taxonomies
 - Bugs models
 - Vulnerability models

BF Weakness

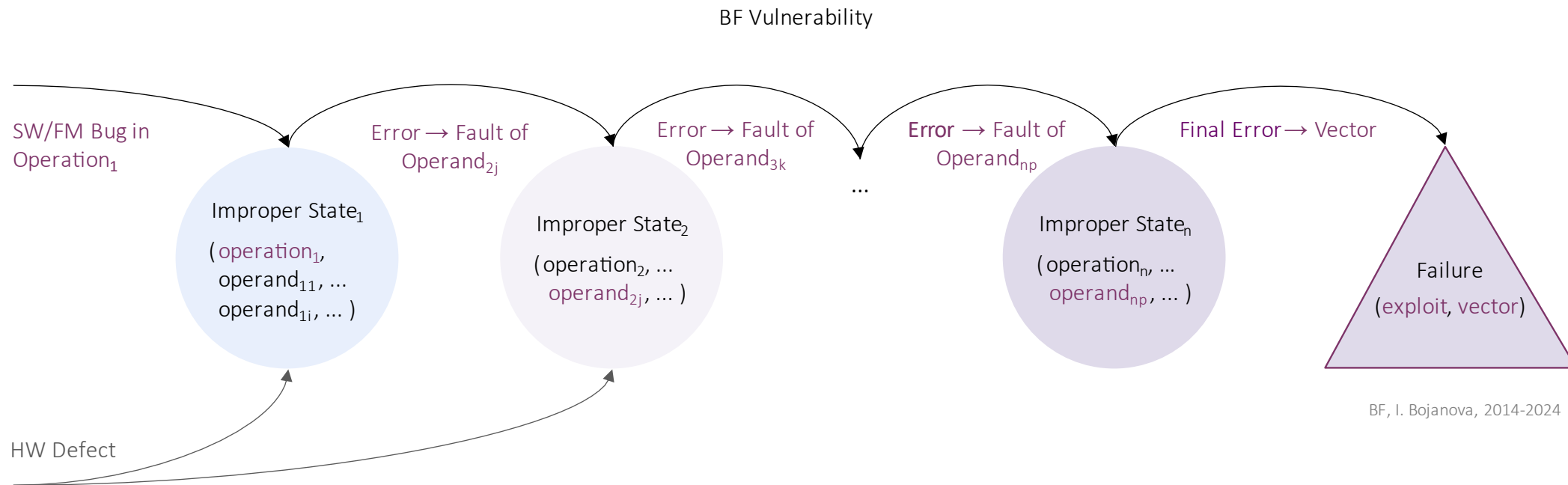


BF Weakness States



-  Improper State caused by a Bug – the operation is improper
-  Improper State caused by a Fault – an operand is improper



BF Vulnerability





BF, I. Bojanova, 2014-2024

Improper State: an $(operation, operand, \dots, operand_n)$ tuple with at least one improper element

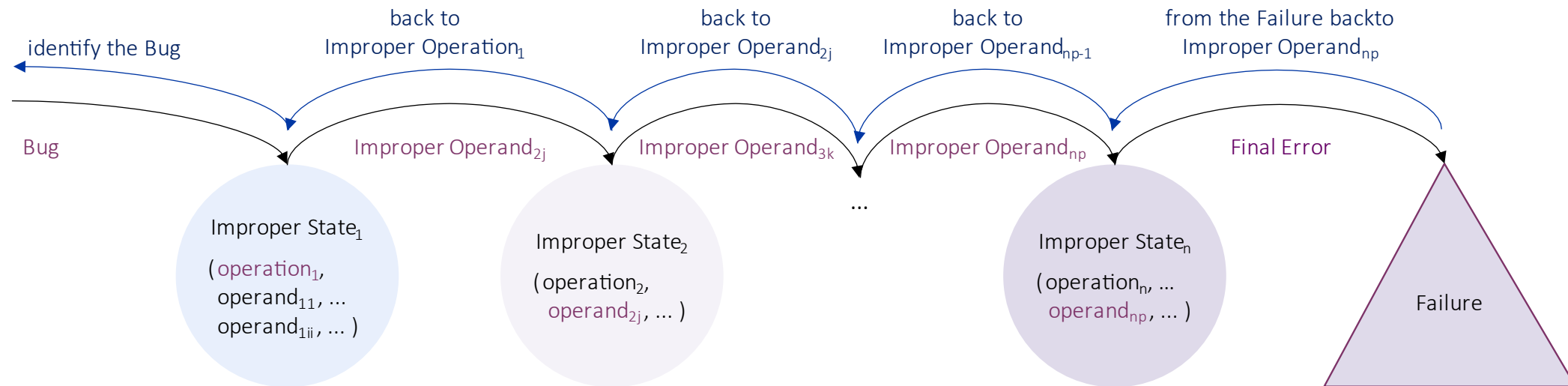
↻ Chaining

-  Initial State – caused by a Bug
-  Propagation State – caused by a Fault

-  Final State – supplies an Exploit Vector
-  Failure – result of the exploit of the vector supplied by the Final Error

BF Bugs Detection

BF Bug Identification



BF, I. Bojanova, 2014-2024

Improper State: an $(operation, operand_1, \dots, operand_n)$ tuple with at least one improper element

↻ Chaining ↶ Backtrack to previous State

- Initial State – caused by a Bug
- Propagation State – caused by a Fault

- Final State – results in an Exploitable Error
- Failure – caused by exploitation of the Final Error

Bug/Fault – relates to Execution Phase:

Operations

Input Operands

Output Results

- **Security Bug**

- Code or specification defect
- May result from a hardware defect
- May resurface by configuration/environment

- **Fault**

- Name, data, type, address, or size error
- Could be from a Bug or induced by a hardware defect

- **Error**

- From bug or fault
- Propagates to another fault

- **Security Final Error**

- From bug or fault
- Undefined system behavior

- **Security Weakness**

- (bug, operation, error)
(fault, operation, error)
(bug, operation, final error)
(fault, operation, final error)

- **Security Vulnerability**

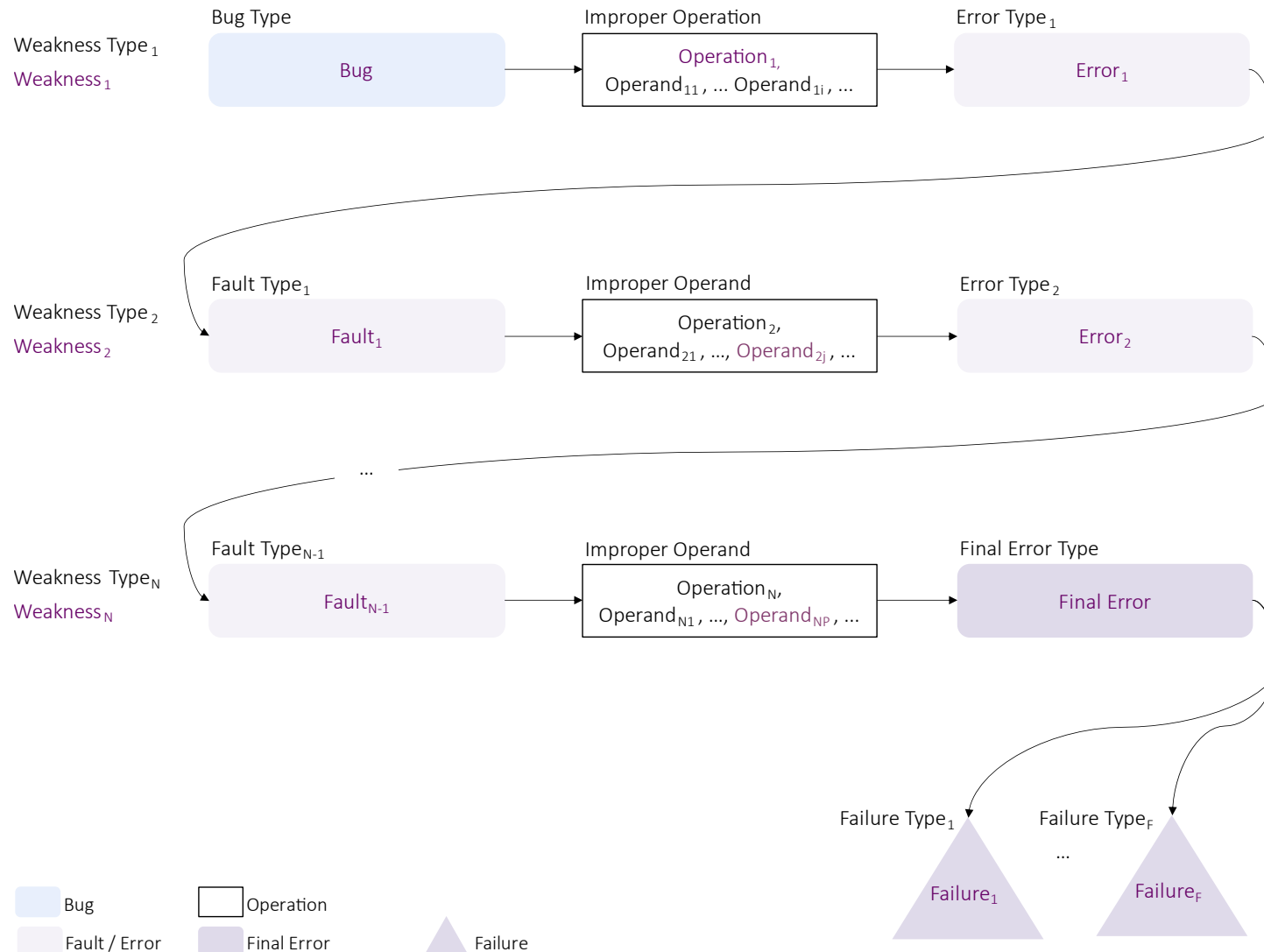
- Chain of weaknesses
- Bug → Error/Fault → ... → Final Error

- **Security Failure**

- Violation of system security requirement
 - Information Exposure (IEX)
 - Data Tempering (TPR)
 - Denial of Service (DoS)
 - Arbitrary Code Execution (ACE)

BF Security Concepts Model

BF Software Security Concepts Model

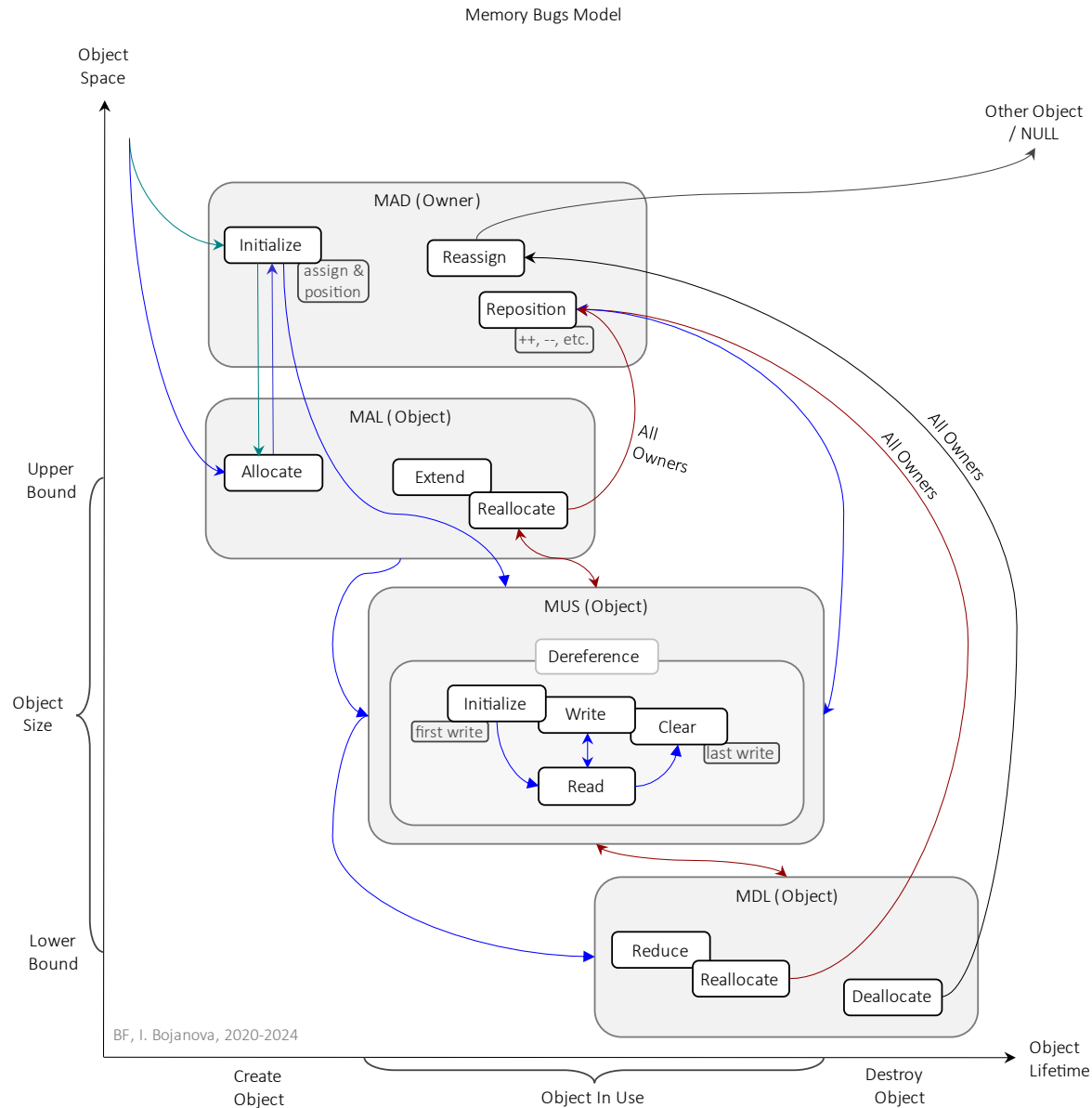


Operation₁ has a Bug and results in Error 1, which becomes Fault₁ for Operation₂, leading to Error₂.

The chain goes on, until the last operation results in a Final Error, leading to a Failure.

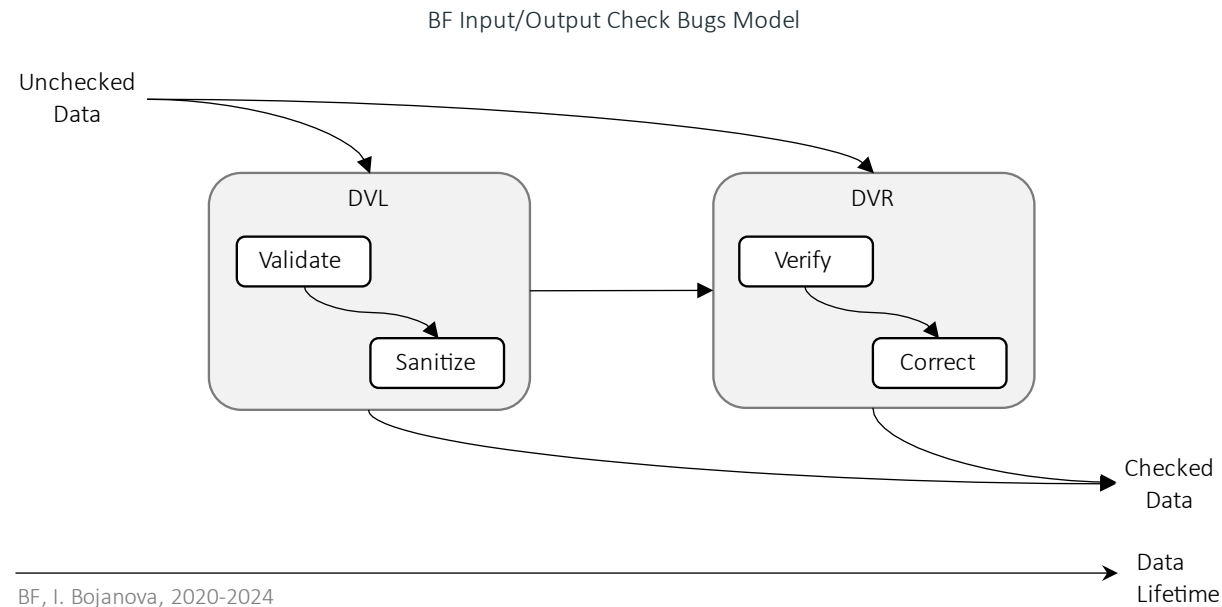
BF Bugs Models

BF Memory (_MEM) Bugs Models



- Identify Secure Code Principles:
 - Memory Safety

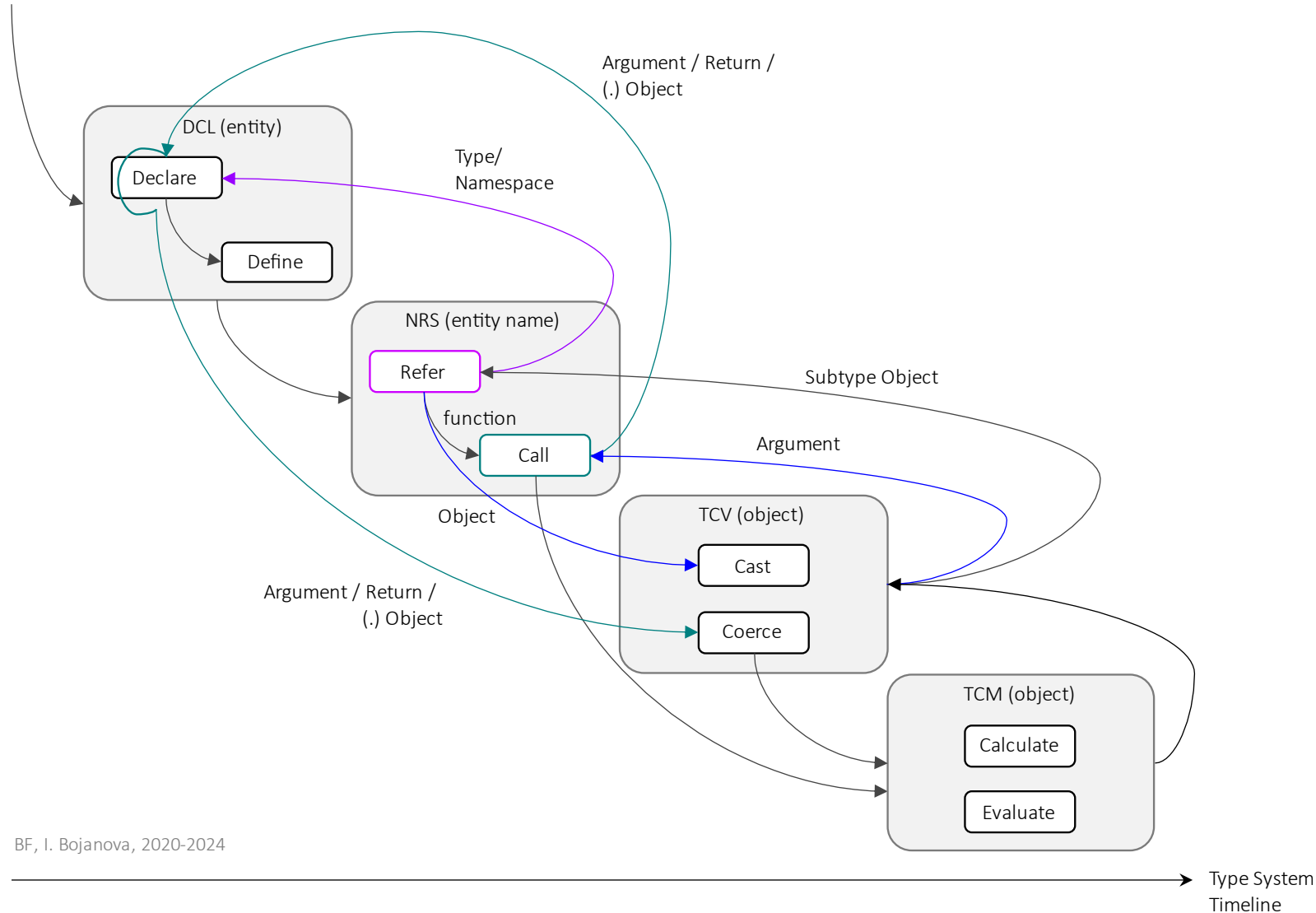
BF Input/Output Model (_INP) Bugs Model



- Identify Secure Code Principles:
 - Input/Output Safety

BF Data Type (_DAT) Bugs Model

BF Data Type Bugs Model



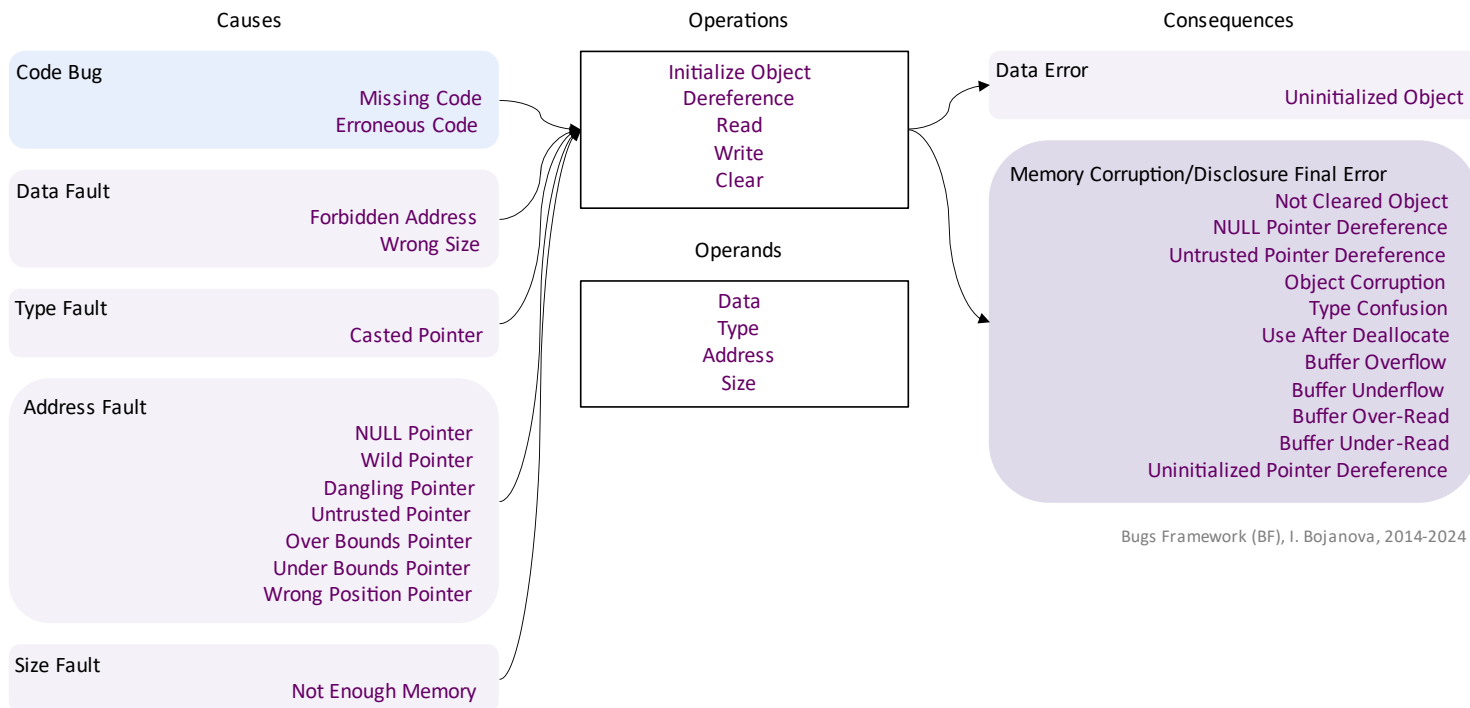
- Identify Secure Code Principles:
 - Data Type Safety

BF Weakness Taxonomies

BF Memory Use (MUS) Class – Example from _MEM Class Type

Bugs Framework (BF)

Memory Corruption/Disclosure (_MEM) Class Type
Memory Use (MUS) Class



Bugs Framework (BF), I. Bojanova, 2014-2024

BF Memory Use (MUS) Class

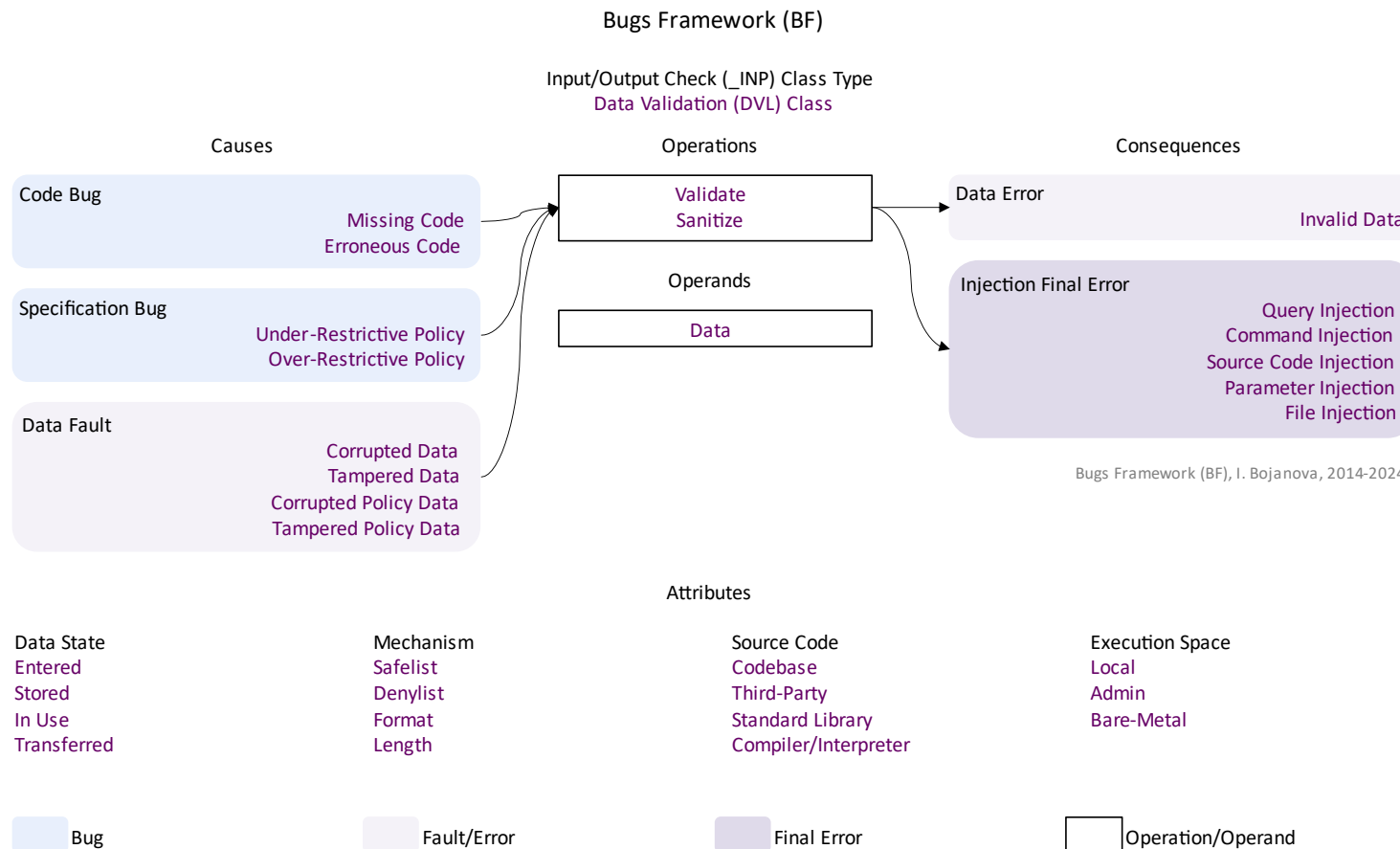
An object is initialized, read, written, or cleared improperly.

- (Bug, Operation₁, Error₁) ← lookup_weakness_triple()
- (Fault₁, Operation₂, Error₂) ← lookup_weakness_triple()
- (Fault_{n-1}, Operation_n, Final Error) ← lookup_weakness_triple()

Address Kind	Address State	Size Kind	Mechanism	Source Code	Execution Space
Huge	Stack	Actual	Direct	Codebase	Userland
Moderate	Heap	Used	Sequential	Third-Party	Kernel
Little	/other/			Standard Library	Bare-Metal
				Compiler/Interpreter	

<https://samate.nist.gov/BF/>Taxonomy>

BF Data Validation (DVL) Class – Example from _INP Class Type

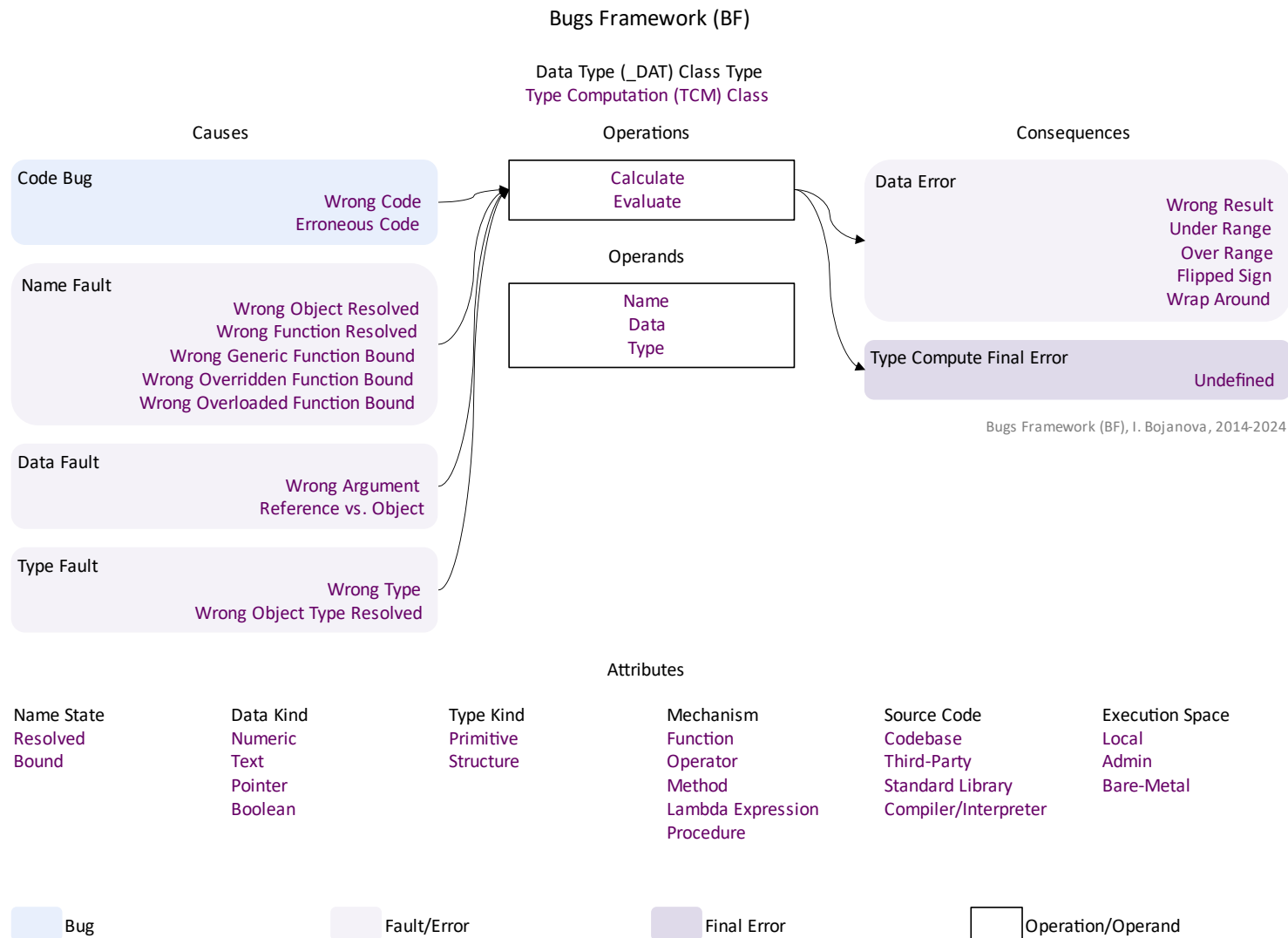


BF Data Validation (DVL) Class

Data are validated (syntax check) or sanitized (escape, filter, repair) improperly.

- (Bug, Operation₁, Error₁) ← lookup_weakness_triple()
- (Fault₁, Operation₂, Error₂) ← lookup_weakness_triple()
- (Fault_{n-1}, Operation_n, Final Error) ← lookup_weakness_triple()

BF Type Computation (TCM) Class – Example from _DAT Class Type



Type Computation (TCM) Class

An arithmetic expression (over numbers, strings, or pointers) is calculated improperly, or a boolean condition is evaluated improperly.

- (Bug, Operation₁, Error₁) ← lookup_weakness_triple()
- (Fault₁, Operation₂, Error₂) ← lookup_weakness_triple()
- (Fault_{n-1}, Operation_n, Final Error) ← lookup_weakness_triple()

BF Weakness Taxonomies

- **Structured**
(bug/fault, operation, error/final error)
- **Complete**
no gaps in coverage
- **Orthogonal**
no overlaps
- **Language and domain independent**
context-free
- **Causation rules**
cause-consequence transition by operation

BF Taxonomy – BF.xml

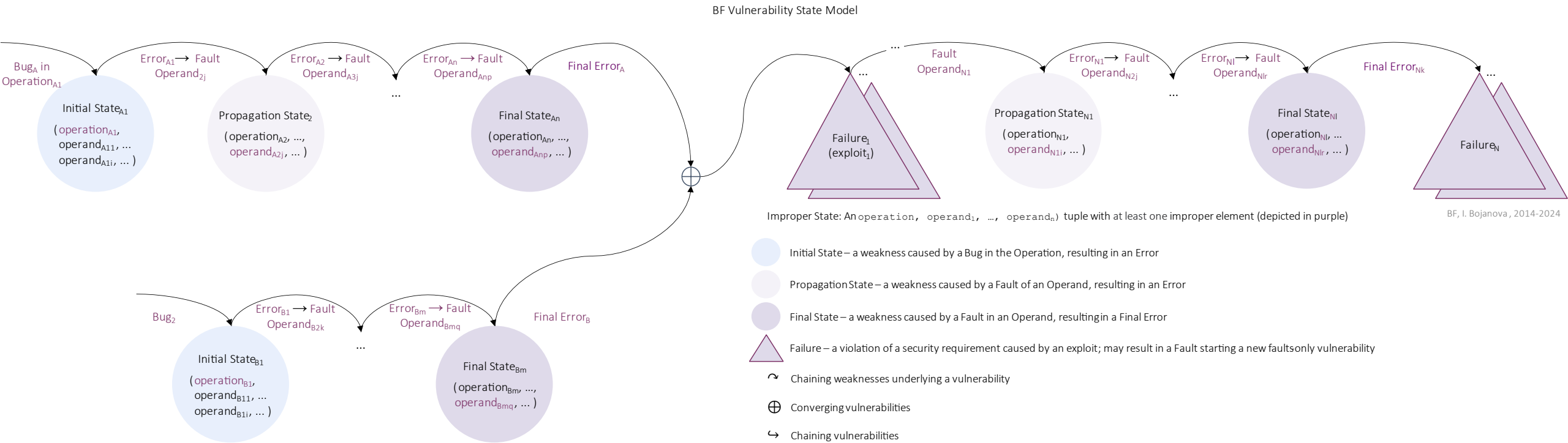
```
BF.xml* [X]
<!--@author Irena Bojanova(ivb)-->
<!--@date - 2/9/2022-->
<BF Name="Bugs Framework">
  <Cluster Name="_INP" Type="Weakness">...</Cluster>
  <Cluster Name="_DAT" Type="Weakness">
    <Class Name="DCL" Title="Declaration Bugs">
      <Operations>
        <Operation Name="Declare"/>
        <Operation Name="Define"/>
        <AttributeType Name="Mechanism">...</AttributeType>
        <AttributeType Name="Source Code">...</AttributeType>
        <AttributeType Name="Entity">...</AttributeType>
      </Operations>
      <Operands>
        <Operand Name="Type"><!--XXX-->
          <AttributeType Name="Type Kind">...</AttributeType>
        </Operand>
      </Operands>
      <Causes>
        <BugCauseType Name="The Bug">
          <Cause Name="Missing Code"/>
          <Cause Name="Wrong Code"/>
          <Cause Name="Erroneous Code"/>
          <Cause Name="Missing Modifier"/>
          <Cause Name="Wrong Modifier"/>
          <Cause Name="Anonymous Scope"/>
          <Cause Name="Wrong Scope"/>
        </BugCauseType>
      </Causes>
      <Consequences>
        <WeaknessConsequenceType Name="Improper Type (_DAT)">
          <Consequence Name="Wrong Type"/>
          <Consequence Name="Incomplete Type"/>
        </WeaknessConsequenceType>
      </Consequences>
    </Class>
  </Cluster>
</BF>
```

```
BF.xml* [X]
<Definitions>
  <!-- Clusters-->
  <Definition Name="_INP" Type="Weakness">Input/Output Check Bugs
  <Definition Name="_DAT" Type="Weakness">Data Type Bugs - lead to
  <Definition Name="_MEM" Type="Weakness">Memory Bugs - lead to M
  <Definition Name="_CRY" Type="Weakness">Cryptographic Store or
  <Definition Name="_RND" Type="Weakness">Random Number Generation
  <Definition Name="_ACC" Type="Weakness">Access Control Bugs - le

  <!-- Classes - xxx update the definitions on BF web-site-->
  <!-- _INP-->
  <Definition Name="DVL">Data are validated (syntax check) or san
  <Definition Name="DVR">Data are verified (semantics check) or co
  <!-- _DAT-->
  <Definition Name="DCL">An object, a function, a type, or a name
  <Definition Name="NRS">The name of an object, a function, or a
  <Definition Name="TCV">Data are converted or coerced into other
  <Definition Name="TCM">A numeric, pointer, or string value is ca
  <!-- _MEM-->
  <Definition Name="MAD">The pointer to an object is initialized,
  <Definition Name="MAL">An object is allocated, extended, or rea
  <Definition Name="MUS">An object is initialized, read, written,
  <Definition Name="MDL">An object is deallocated, reduced, or rea
  ...
  <!-- Values-->
  ...
```

BF Vulnerability Models

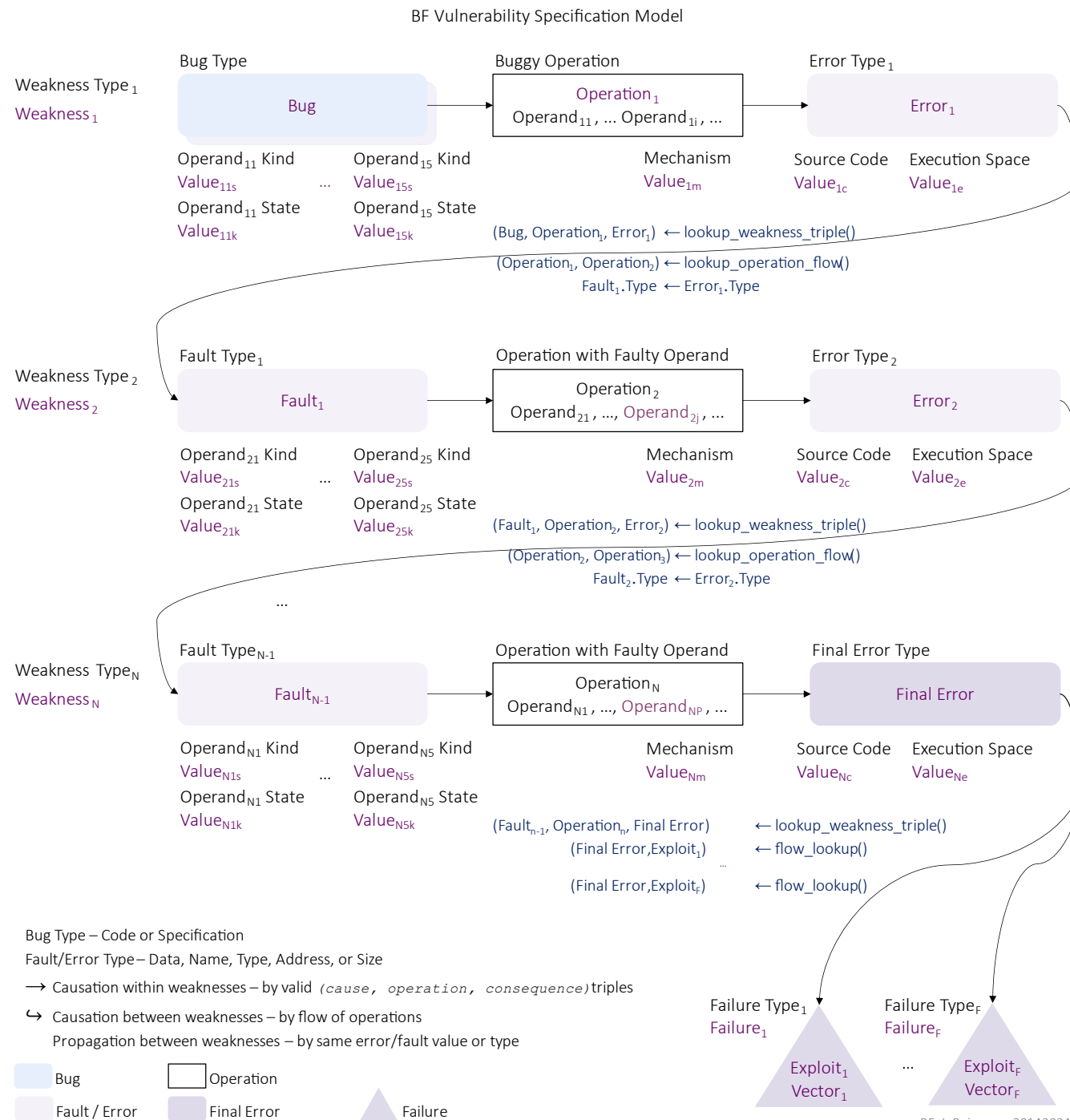
BF Vulnerability State Model



- The bug in at least one of the chains must be fixed to resolve the vulnerability
- Fixing a fault may only mitigate the vulnerability

BF Vulnerability Specification Model

- Chain of (cause, operation, consequence) weakness triples
- Bug = improper operation
- Fault = improper operand
- Bug Types: Code, Specification
- Fault Types: Name, Data, Type, Address, Size
- Causation within a weakness
- Causation between weaknesses
- Causation between vulnerabilities
- Propagation between weaknesses
- Propagation between vulnerabilities



BF Formal Language

BF Context Free Grammar (CFG)

$$G = (V, \Sigma, R, S) \quad (1)$$

, where:

- Σ defines the BF lexis (the alphabet of the CFG) as a finite set of tokens (terminals) comprised by the sets of BF taxons and BF symbols (see Listing 3)

$$\Sigma = \{ \alpha \mid \alpha \in \Sigma Taxon \cup \Sigma Symbol \}$$

- V and R define the BF syntax as
 - a finite set of variables (nonterminals)

$$V = \{ S, V_1, \dots, V_n \}$$

and

- a finite set of syntactic rules (productions) in the form

$$R = \{ A \mapsto \omega \mid A \in V \wedge \omega \in (V \cup \Sigma)^* \}$$

, where:

$(V \cup \Sigma)^*$ is a string of tokens and/or variables

$A \mapsto \omega$ means any variable A occurrence may be replaced by ω .

- $S \in V$ is the predefined start variable from which all BF specifications derive.

The formal language is defined as the set of all strings of tokens ω derivable from the start variable S .

$$L(G) = \{\omega \in \Sigma^* : S \xRightarrow{*} \omega\} \quad (2)$$

, where:

- Σ^* is the set of all possible strings that can be generated from Σ tokens
- S is the start variable
- $\alpha \xRightarrow{*} \beta$ means string α derives string β

Note that ω must be in Σ^* , the set of strings made from terminals. Strings involving non-terminals are not part of the language.

BF CFG Lexis

$$\Sigma = \{\Sigma\text{Taxon}, \Sigma\text{Symbol}\}$$

, where

$$\Sigma\text{Taxon} = \{\Sigma\text{Category}, \Sigma\text{ClassType}, \Sigma\text{Class}, \Sigma\text{BugType}, \Sigma\text{Bug}, \Sigma\text{Operation}, \Sigma\text{OperationAttributeType}, \Sigma\text{FaultType}, \Sigma\text{Fault}, \Sigma\text{OperandAttributeType}, \Sigma\text{OperandAttribute}, \Sigma\text{FinalErrorType}, \Sigma\text{FinalError}\}$$

$$\Sigma\text{Symbol} = \{\rightarrow, \leftrightarrow, \oplus\}$$

$$\Sigma\text{Category} = \{\text{'Weakness'}, \text{'Failure'}\}$$

$$\Sigma\text{ClassType} = \{\text{'_INP'}, \text{'_DAT'}, \text{'_MEM'}, \dots\}$$

$$\Sigma\text{Class} = \{\text{'DVL'}, \text{'DVR'}, \text{'DCL'}, \text{'NRS'}, \text{'TCV'}, \text{'TCM'}, \text{'MAD'}, \text{'MMN'}, \text{'MUS'}, \dots\}$$

$$\Sigma\text{Operation} = \{\text{'Validate'}, \text{'Sanitize'}, \text{'Verify'}, \text{'Correct'}, \text{'Declare'}, \text{'Define'}, \text{'Refer'}, \text{'Call'}, \text{'Cast'}, \text{'Coerce'}, \text{'Calculate'}, \text{'Evaluate'}, \text{'InitializePointer'}, \text{'Reposition'}, \text{'Reassign'}, \text{'Allocate'}, \text{'Extend'}, \text{'Reallocate - Extend'}, \text{'Deallocate'}, \text{'Reduce'}, \text{'Reallocate - Reduce'}, \text{'InitializeObject'}, \text{'Dereference'}, \text{'Read'}, \text{'Write'}, \text{'Clear'}, \text{'Generate/Select'}, \text{'Store'}, \text{'Distribute'}, \text{'Use'...}\}$$

$$(3) \text{Operation} = \{\text{'Validate'}, \text{'Sanitize'}, \text{'Verify'}, \text{'Correct'}, \text{'Declare'}, \text{'Define'}, \text{'Refer'}, \text{'Call'}, \text{'Cast'}, \text{'Coerce'}, \text{'Calculate'}, \text{'Evaluate'}, \text{'InitializePointer'}, \text{'Reposition'}, \text{'Reassign'}, \text{'Allocate'}, \text{'Extend'}, \text{'Reallocate - Extend'}, \text{'Deallocate'}, \text{'Reduce'}, \text{'Reallocate - Reduce'}, \text{'InitializeObject'}, \text{'Dereference'}, \text{'Read'}, \text{'Write'}, \text{'Clear'}, \text{'Generate/Select'}, \text{'Store'}, \text{'Distribute'}, \text{'Use'...}\}$$

$$\text{BugType} = \{\text{'CodeDefect'}, \text{'SpecificationDefect'}\}$$

$$\Sigma\text{Bug} = \{\text{'MissingCode'}, \text{'ErroneousCode'}, \text{'Under - RestrictivePolicy'}, \text{'Over - RestrictivePolicy'}, \text{'WrongCode'}, \text{'MissingModifier'}, \text{'WrongModifier'}, \text{'AnonymousScope'}, \text{'WrongScope'}, \text{'MissingQualifier'}, \text{'WrongQualifier'}, \text{'MismatchedOperation'}, \dots\}$$

$$\Sigma\text{FinalErrorType} = \{\text{'Injection'}, \text{'Access'}, \text{'TypeCompute'}, \text{'MemoryCorruption/Disclosure'}, \dots\}$$

$$\Sigma\text{FinalError} = \{\text{'QueryInjection'}, \text{'CommandInjection'}, \text{'SourceCodeInjection'}, \text{'ParameterInjection'}, \text{'FileInjection'}, \text{'WrongAccessObject'}, \text{'WrongAccessType'}, \text{'WrongAccessFunction'}, \text{'Undefined'}, \text{'MemoryLeak'}, \text{'MemoryOverflow'}, \text{'DoubleDeallocate'}, \text{'ObjectCorruption'}, \text{'NotClearedObject'}, \text{'NULLPointerDereference'}, \text{'UntrustedPointerDereference'}, \text{'TypeConfusion'}, \text{'UseAfterDeallocate'}, \text{'BufferOverflow'}, \text{'BufferUnderflow'}, \text{'BufferOver - Read'}, \text{'BufferUnder - Read'}\}$$

BF CFG Syntax

$$S ::= (Vulnerability (\oplus Vulnerability)? Failure+) + \epsilon \quad (4)$$

Vulnerability ::= + *Weakness*

Weakness ::= *Cause Operation Consequence*

Cause ::= *Bug* | *Fault*

Consequence ::= *Error* | *FinalError*

$$S ::= (Vulnerability (\oplus Vulnerability)? Failure+) + \epsilon \quad (5)$$

Vulnerability ::= *SingleWeakness*

| *FirstWeakness* (*Weakness*+) *LastWeakness*

SingleWeakness ::= *Bug Operation FinalError*

FirstWeakness ::= *Bug Operation* (*Error* | *FinalError*)

Weakness ::= *Fault Operation Error*

LastWeakness ::= *Error Operation FinalError*

BF Syntax – LL(1) Grammar

$S ::= Vulnerability Converge_Failure$ (6)

$Vulnerability ::= Bug_Fault Operation OperAttrs_Error_FinalError$

$Bug_Fault ::= Bug$

$| Fault$

$OperAttrs_Error_FinalError ::= OperationAttribute OperAttrs_Error_FinalError$

$| Error Fault OprndAttrs_Operation$

$| FinalError$

$OprndAttrs_Operation ::= OperandAttribute OprndAttrs_Operation$

$| Operation OperAttrs_Error_FinalError$

$Converge_Failure ::= \oplus Vulnerability Converge_Failure$

$| Vector Exploit NextVulner_Failure$

$NextVulner_Failure ::= Fault OprndAttrs_Operation$

$| Failure \epsilon$

BF Semantics – Attribute CGF

SyntaxRules :

$S ::= \text{Vulnerability Converge_Failure}$

$\text{Vulnerability} ::= \text{Bug_Fault Operation OperAttrs_Error_FinalError}$

$\text{Bug_Fault} ::= \text{Bug}$

| Fault

$\text{OperAttrs_Error_FinalError} ::= \text{OperationAttribute OperAttrs_Error_FinalError}$

| $\text{Error Fault}_1 \text{OprndAttrs_Operation}$

| FinalError

$\text{OprndAttrs_Operation} ::= \text{OperandAttribute OprndAttrs_Operation}$

| $\text{Operation}_k \text{OperAttrs_Error_FinalError}$

$\text{Converge_Failure} ::= \oplus \text{Vulnerability Converge_Failure}$

| $\text{Vector Exploit NextVulner_Failure}$

$\text{NextVulner_Failure} ::= \text{Fault}_2 \text{OprndAttrs_Operation}$

| $\text{Failure } \varepsilon$

SemanticRules :

$(\text{Bug}, \text{Operation}_1, \text{Error}) \leftarrow \text{lookup_weakness_triple}()$

$(\text{Bug}, \text{Operation}_1, \text{FinalError}) \leftarrow \text{lookup_weakness_triple}()$

$(\text{Fault}_1, \text{Operation}_k, \text{Error}), k > 1 \leftarrow \text{lookup_weakness_triple}()$

$(\text{Fault}_1, \text{Operation}_k, \text{FinalError}), k > 1 \leftarrow \text{lookup_weakness_triple}()$

$(\text{Operation}_1, \dots, \text{Operation}_k), k > 1 \leftarrow \text{lookup_operation_flow}()$

$\text{Fault}_1 \leftarrow \text{if } (\text{Fault}_1.\text{ClassType} == \text{Error}.\text{ClassType}) \text{ then Error}$

Predicates :

$\text{Fault}_1.\text{Type} == \text{Error}.\text{Type}$

$\text{Vector}.\text{Type} == \text{FinalError}.\text{Type}$

$\text{Fault}_2.\text{Type} == \text{ExploitResult}.\text{Type}$

(7)

BF Specifications of CWEs

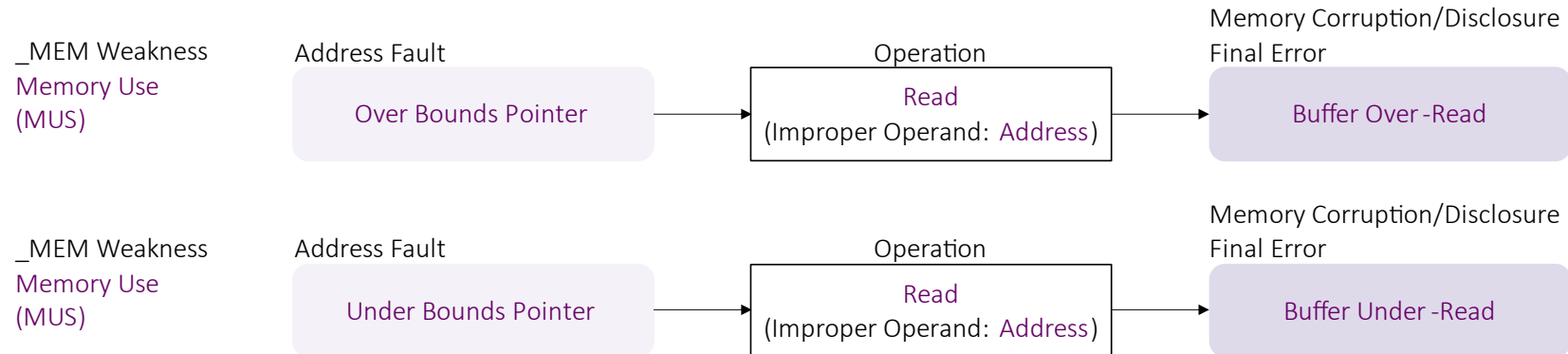
BFCWE Dataset

```
<?xml version="1.0" encoding="utf-8"?>
<!--Bugs Framework (BF), BFCWE Tool, I. Bojanova, NIST, 2020-2024-->
<BFCWE-Dataset>
  <CWE ID="20">
    <BFCWE Cause="Missing Code" Operation="Verify" Consequence="Wrong Value" />
    <BFCWE Cause="Missing Code" Operation="Verify" Consequence="Inconsistent Value" />
    <BFCWE Cause="Missing Code" Operation="Verify" Consequence="Wrong Type" />
    <BFCWE Cause="Erroneous Code" Operation="Verify" Consequence="Wrong Value" />
    <BFCWE Cause="Erroneous Code" Operation="Verify" Consequence="Inconsistent Value" />
    <BFCWE Cause="Erroneous Code" Operation="Verify" Consequence="Wrong Type" />
    <BFCWE Cause="Missing Code" Operation="Validate" Consequence="Invalid Data" />
    <BFCWE Cause="Erroneous Code" Operation="Validate" Consequence="Invalid Data" />
  </CWE>
  <CWE ID="22">
    <BFCWE Cause="Under-Restrictive Policy" Operation="Sanitize" Consequence="File Injection" />
  </CWE>
  <CWE ID="23">
    <BFCWE Cause="Under-Restrictive Policy" Operation="Sanitize" Consequence="File Injection" />
  </CWE>
  <CWE ID="24">
    <BFCWE Cause="Under-Restrictive Policy" Operation="Sanitize" Consequence="File Injection" />
  </CWE>
  <CWE ID="25">
    <BFCWE Cause="Under-Restrictive Policy" Operation="Sanitize" Consequence="File Injection" />
  </CWE>
  <CWE ID="26">
    <BFCWE Cause="Under-Restrictive Policy" Operation="Sanitize" Consequence="File Injection" />
  </CWE>
  <CWE ID="27">
    <BFCWE Cause="Under-Restrictive Policy" Operation="Sanitize" Consequence="File Injection" />
  </CWE>
  <CWE ID="28">
    <BFCWE Cause="Under-Restrictive Policy" Operation="Sanitize" Consequence="File Injection" />
  </CWE>
  <CWE ID="29">
    <BFCWE Cause="Under-Restrictive Policy" Operation="Sanitize" Consequence="File Injection" />
  </CWE>
  <CWE ID="30">
    <BFCWE Cause="Under-Restrictive Policy" Operation="Sanitize" Consequence="File Injection" />
  </CWE>
  <CWE ID="31">
```

```
<CWE ID="125">
  <BFCWE Cause="Over Bounds Pointer" Operation="Read" Consequence="Buffer Over-Read" />
  <BFCWE Cause="Under Bounds Pointer" Operation="Read" Consequence="Buffer Under-Read" />
</CWE>
<CWE ID="126">
  <BFCWE Cause="Over Bounds Pointer" Operation="Read" Consequence="Buffer Over-Read" />
</CWE>
<CWE ID="127">
  <BFCWE Cause="Under Bounds Pointer" Operation="Read" Consequence="Buffer Under-Read" />
</CWE>
<CWE ID="128">
  <BFCWE Operation="Calculate" Consequence="Wrap Around" />
</CWE>
<CWE ID="129">
  <BFCWE Cause="Missing Code" Operation="Verify" Consequence="Wrong Value" />
  <BFCWE Cause="Erroneous Code" Operation="Verify" Consequence="Wrong Value" />
</CWE>
<CWE ID="130">
  <BFCWE Cause="Missing Code" Operation="Verify" Consequence="Inconsistent Value" />
  <BFCWE Cause="Erroneous Code" Operation="Verify" Consequence="Inconsistent Value" />
</CWE>
<CWE ID="131">
  <BFCWE Operation="Calculate" />
</CWE>
<CWE ID="134">
  <BFCWE Cause="Missing Code" Operation="Validate" Consequence="Parameter Injection" />
</CWE>
<CWE ID="135">
  <BFCWE Operation="Calculate" Consequence="Wrong Result" />
</CWE>
<CWE ID="138">
  <BFCWE Cause="Missing Code" Operation="Sanitize" Consequence="Invalid Data" />
  <BFCWE Cause="Erroneous Code" Operation="Sanitize" Consequence="Invalid Data" />
</CWE>
<CWE ID="140">
  <BFCWE Cause="Missing Code" Operation="Sanitize" Consequence="Parameter Injection" />
  <BFCWE Cause="Erroneous Code" Operation="Sanitize" Consequence="Parameter Injection" />
</CWE>
```

CWE-125 – Two BF Specifications

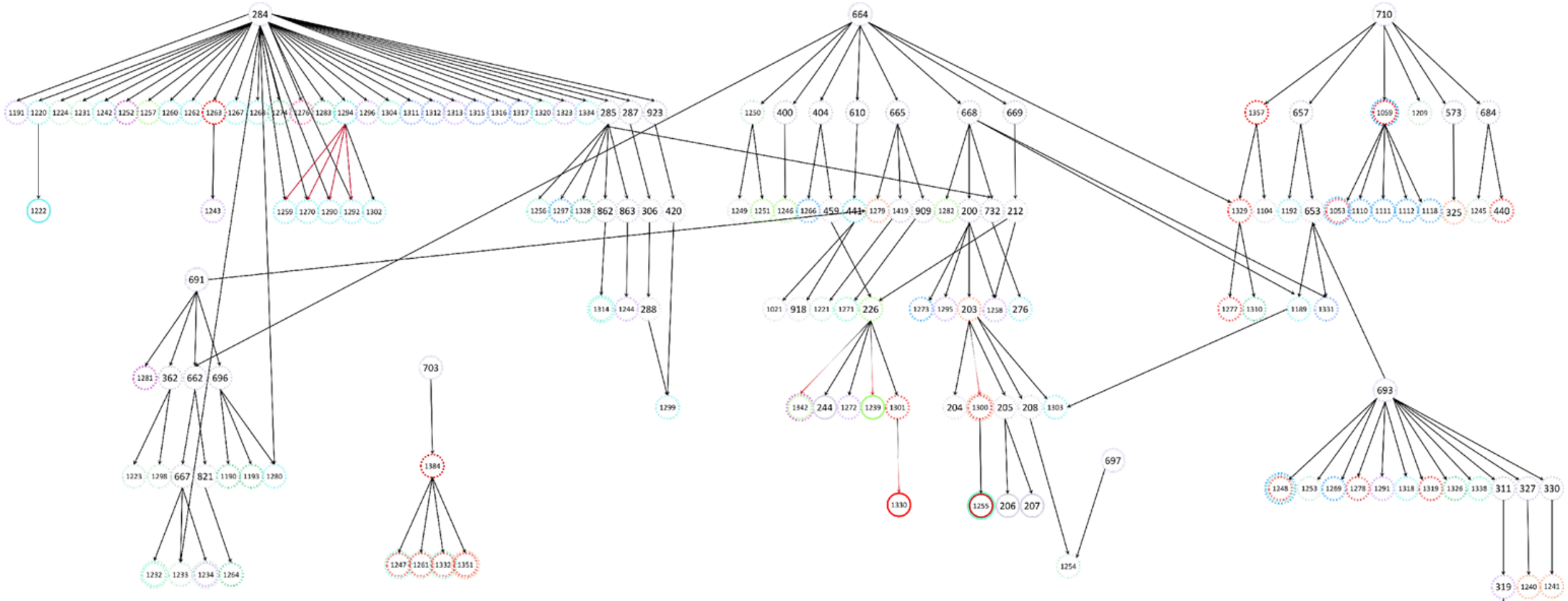
BF Specifications of CWE-125



BF, I. Bojanova, 2014-2024

Class Type	Definition
Memory Corruption/Disclosure (_MEM)	Bugs/weaknesses allowing a memory corruption/disclosure exploit.
Class	Definition
Memory Use (MUS)	An object is initialized, read, written, or cleared improperly.
Operation	Definition
Read	Use the value of an object's data.
Cause	Definition
Address Fault	The object address in use is wrong.
Over Bounds Pointer	Holds an address above the upper boundary of its object.
Under Bounds Pointer	Holds an address above the lower boundary of its object.
Consequence	Definition
Memory Corruption/Disclosure Final Error	An exploitable or undefined system behavior caused by memory addressing, allocation, use, and deallocation bugs.
Buffer Over-Read	Reading above the upper bound of an object.
Buffer Under-Read	Reading below the lower bound of an object.

Analyzing HW CWEs



CWEs from:

- Manufacturing and Life Cycle Management Concerns - (1195)
- Security Flow Issues - (1196)
- Integration Issues - (1197)
- Privilege Separation and Access Control Issues - (1198)

○

- General Circuit and Logic Design Concerns - (1199)
- Core and Compute Issues - (1201)
- Memory and Storage Issues - (1202)
- Peripherals, On-chip Fabric, and Interface/IO Problems - (1203)
- Security Primitives and Cryptography Issues - (1205)

○

- Power, Clock, Thermal, and Reset Concerns - (1206)
- Debug and Test Problems - (1207)
- Cross-Cutting Problems - (1208)
- Physical Access Issues and Concerns - (1388)
- Not Hardware CWEs

CWEs from:

- / Only View-1000
- / Only View-1194
- / Both View-1000 and View-1194

CWEs by Abstraction:

- Pillar
- Class
- Base
- Variant

CWEs Relation by Nature:

- / ChildOf

BF Specifications of CVEs

Heartbleed (CVE-2014-0160)

[CVE-2014-0160](#)

The (1) TLS and (2) DTLS implementations in OpenSSL 1.0.1 before 1.0.1g do not properly handle Heartbeat Extension packets, which allows remote attackers to obtain sensitive information from process memory via crafted packets that trigger a **buffer over-read**, as demonstrated by **reading private keys**, related to d1_both.c and t1_lib.c, aka the Heartbleed bug.



→ ↻ <https://nvd.nist.gov/vuln/detail/CVE-2014-0160>

Weakness Enumeration

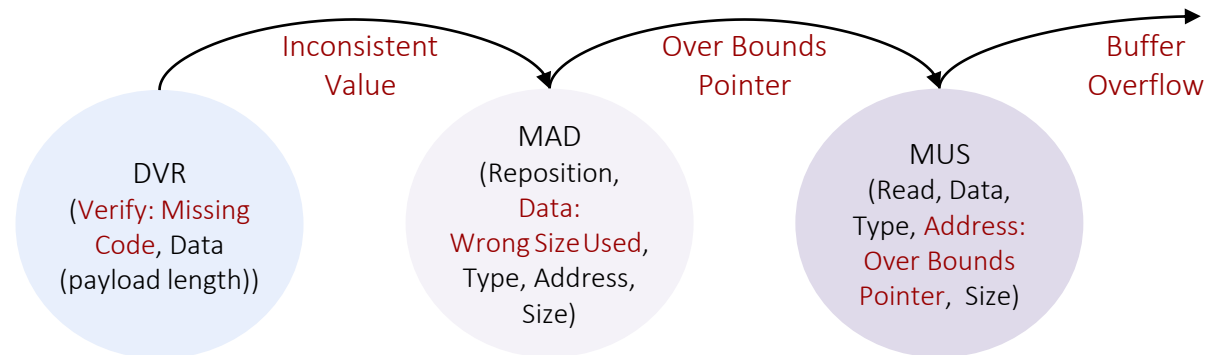
CWE-ID	CWE Name
CWE-119	Improper Restriction of Operations within the Bounds of a Memory Buffer

Heartbleed (CVE-2014-0160)

[CVE-2014-0160](#) The (1) TLS and (2) DTLS implementations in OpenSSL 1.0.1 before 1.0.1g do not properly handle Heartbeat Extension packets, which allows remote attackers to obtain sensitive information from process memory via crafted packets that trigger a **buffer over-read**, as demonstrated by **reading private keys**, related to d1_both.c and t1_lib.c, aka the Heartbleed bug.

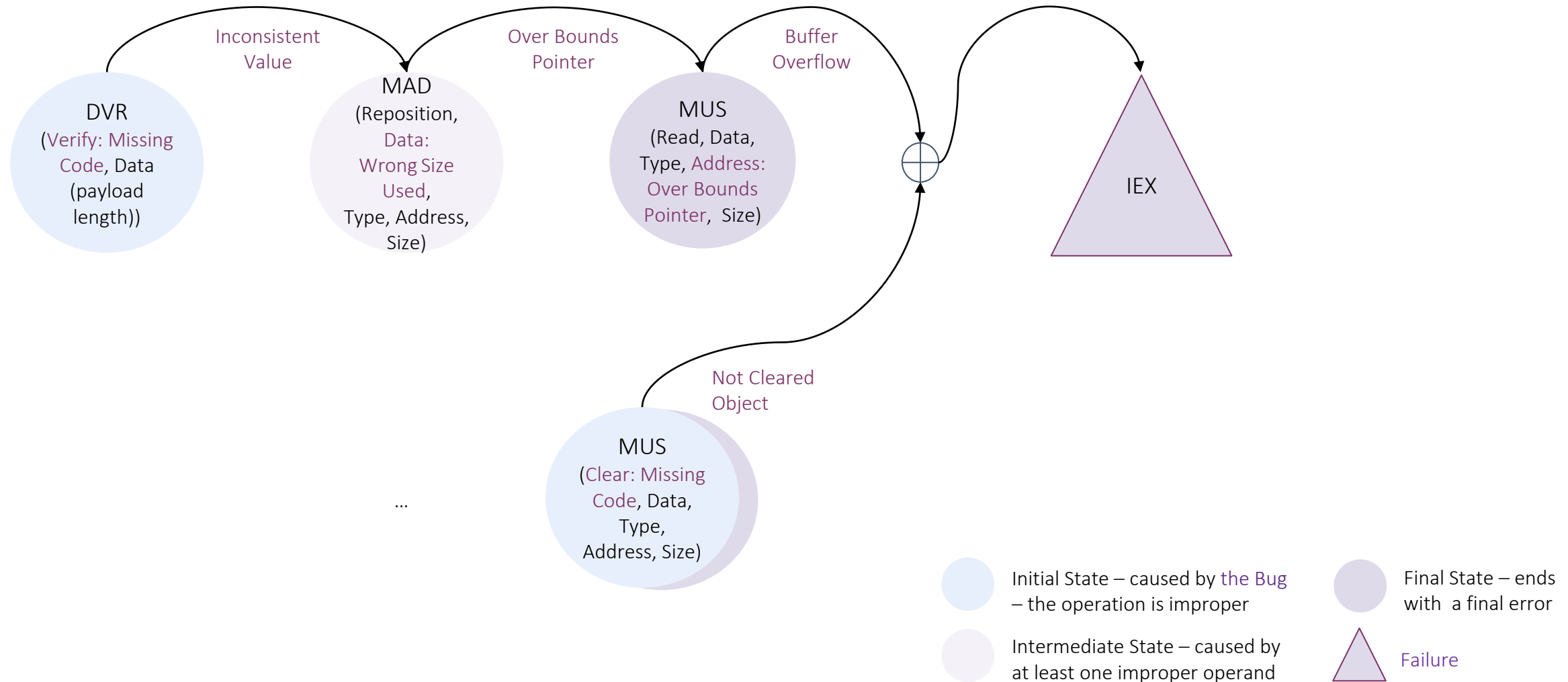
```
1448 dtls1_process_heartbeat(SSL *s)
1449     {unsigned char *p = &s->s3->rrec.data[0], *pl;
1451     unsigned short hbtype;
1452     unsigned
1453     int payload;
1454     unsigned int padding = 16; /* Use minimum padding */
1455
1456     /* Read type and payload length first */
1457     hbtype = *p++;
1458     n2s(p, payload);
1459     pl = p;
1460
1461     ...
1462     if (hbtype == TLS1_HB_REQUEST)
1463     {
1464         unsigned char *buffer, *bp;
1465
1466         ...
1467         /* Allocate memory for the response, size is 1 byte
1468         * message type, plus 2 bytes payload, plus
1469         * payload, plus padding
1470         */
1471         buffer = OPENSSL_malloc(1 + 2 + payload + padding);
1472         bp = buffer;
1473
1474         /* Enter response type, length and copy payload */
1475         *bp++ = TLS1_HB_RESPONSE;
1476         s2n(payload, bp);
1477         memcpy(bp, pl, payload);
1478     }
1479
1480     ...
```

```
/* Naive implementation of memcpy
void *memcpy (void *dst, const void *src, size_t n)
{
    size_t i;
    for (i=0; i<n; i++)
        *(char *) dst++ = *(char *) src++;
    return dst;
}
```

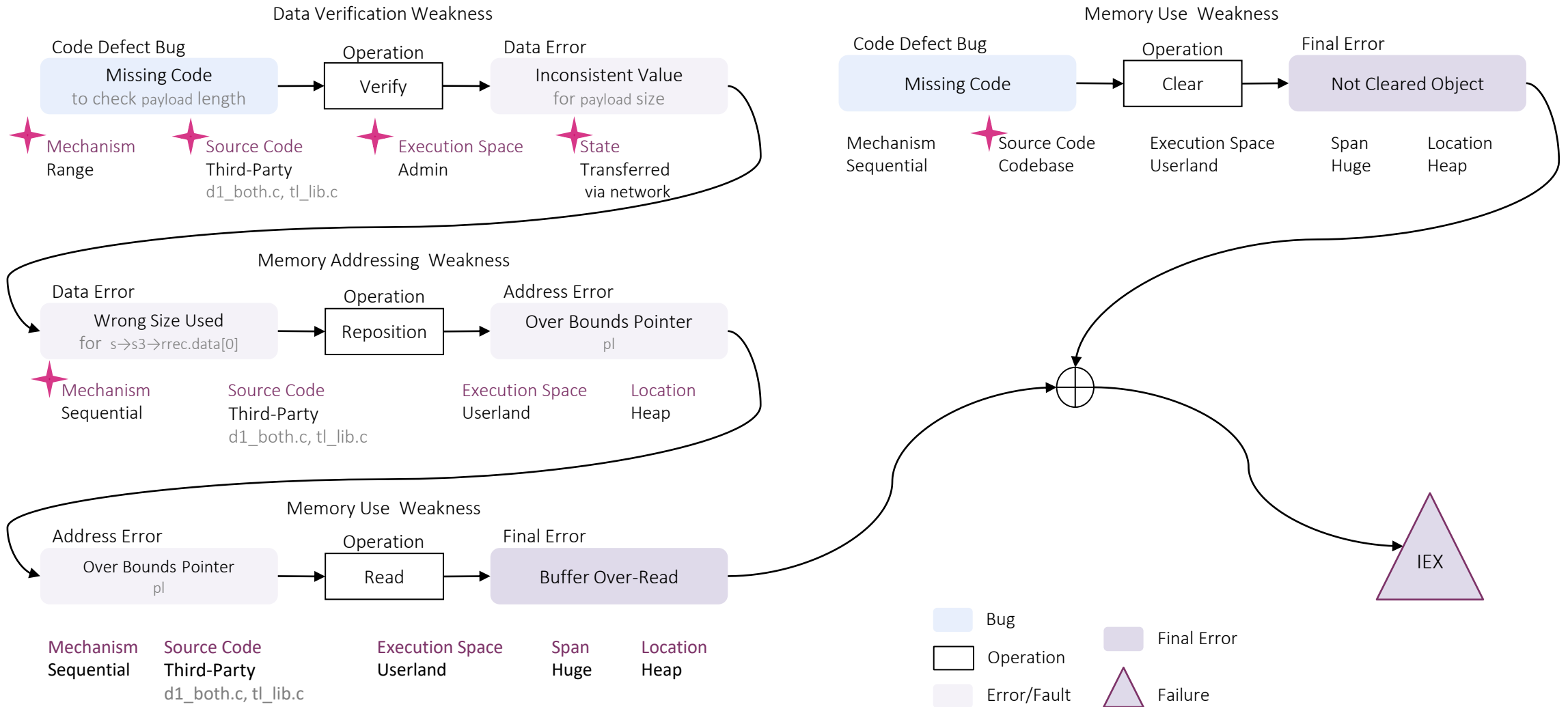


- Caused by the Bug
- Caused by an improper operand
- Ends with a final error

BF States of CVE-2014-0160 (Heartbleed)



BF Specification of CVE-2014-0160 (Heartbleed)



BF Tool – BF Specification of Heartbleed

BF Specification of CVE-2014-0160

BF Title: Heartbleed Heap Buffer Over-Read in OpenSSL v1.0.1 before v1.0.1g

BF Description: Missing verification of 'payload' towards an upper limit leads to use of an inconsistent size for an object, allowing a pointer repetition over its bounds, which, when used in 'memcpy()' leads to a heap buffer over-read. If exploited, this can lead to exposure of sensitive information - confidentiality loss.

Weakness: Over Bounds Pointer

Operation: Instance Pointer

Error / Final Exploitable Error: Wrong Value

BF Specification of CVE-2014-0160

BF Title: Heartbleed Heap Buffer Over-Read in OpenSSL v1.0.1 before v1.0.1g

BF Description: Missing verification of 'payload' towards an upper limit leads to use of an inconsistent size for an object, allowing a pointer repetition over its bounds, which, when used in 'memcpy()' leads to a heap buffer over-read. If exploited, this can lead to exposure of sensitive information - confidentiality loss.

Weakness: Over Bounds Pointer

Operation: Instance Pointer

Error / Final Exploitable Error: Wrong Value

BF Specification of CVE-2014-0160

BF Title: Heartbleed Heap Buffer Over-Read in OpenSSL v1.0.1 before v1.0.1g

BF Description: Missing verification of 'payload' towards an upper limit leads to use of an inconsistent size for an object, allowing a pointer repetition over its bounds, which, when used in 'memcpy()' leads to a heap buffer over-read. If exploited, this can lead to exposure of sensitive information - confidentiality loss.

Weakness: Over Bounds Pointer

Operation: Instance Pointer

Error / Final Exploitable Error: Wrong Value

CVE-2014-0160 - Heartbleed.bfcve

```
CVE-2014-016...rtbleed.bfcve x
<?xml version="1.0" encoding="utf-8"?>
<CVE Name="1 CVE-2014-0160">
  <BugWeakness Type="_INP" Class="DVR">
    <Cause Type="The Bug">Missing Code</Cause>
    <Operation>Verify</Operation>
    <Consequence Comment="for payload size" Type="Improper Data">Inconsistent Value</Consequence>
    <Attributes>...</Attributes>
  </BugWeakness>
  <Weakness Type="_MEM" Class="MAD">
    <Cause Comment="(for s>s3>rrec.data[0])" Type="Improper Data">Wrong Size Used</Cause>
    <Operation>Reposition</Operation>
    <Consequence Type="Improper Address">Over Bounds Pointer</Consequence>
    <Attributes>
      <Operation>
        <Attribute Type="Mechanism">Sequential</Attribute>
        <Attribute Comment="dl_both.c and tl_lib.c" Type="Source Code">Codebase</Attribute>
        <Attribute Type="Execution Space">Userland</Attribute>
      </Operation>
      <Operand Name="Object Address">
        <Attribute Type="Location">Heap</Attribute>
      </Operand>
    </Attributes>
  </Weakness>
  <Weakness Type="_MEM" Class="MUS">
    <Cause Comment="(for s>s3>rrec.data[0])" Type="Improper Address">Over Bounds Pointer</Cause>
    <Operation>Read</Operation>
    <Consequence Type="Memory Error">Buffer Overflow</Consequence>
    <Attributes>...</Attributes>
  </Weakness>
  <Failure Type="_FLR" Class="IEX">
    <Cause Type="Memory Error">Buffer Overflow</Cause>
```

> TAXONOMY

▾ BF CVE

Overview

CVE-2004-1287

CVE-2006-2362

CVE-2007-1320

CVE-2007-6429

CVE-2008-4539

CVE-2013-4930

CVE-2013-4934

CVE-2014-0160

CVE-2015-0235

CVE-2015-5221

CVE-2017-17833

CVE-2018-14557

CVE-2019-14814

CVE-2021-21834

CVE-2022-34835

CVE-2023-1283

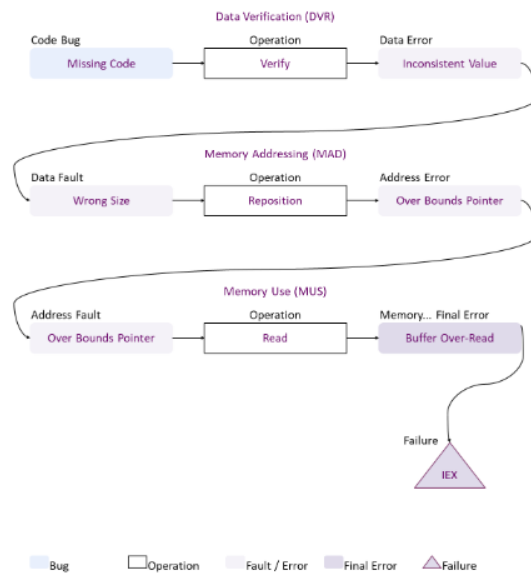
CVE-2023-2356

CVE-2023-2564

CVE-2023-3765

...

BF Specification of CVE-2014-0160 Heartbleed Heap Buffer Over-Read in OpenSSL v1.0.1 before v1.0.1g



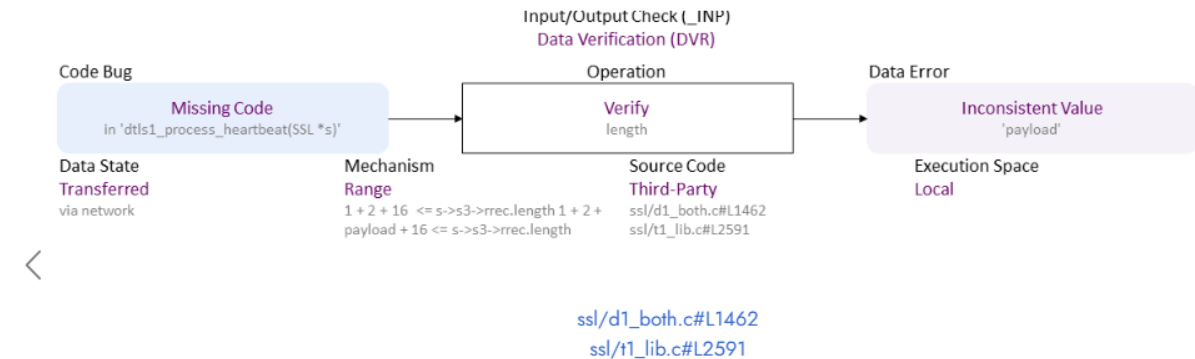
Missing verification of 'payload' towards a upper limit leads to use of an inconsistent size for an object, allowing a pointer reposition over its bounds, which, when used in 'memcpy()' leads to a heap buffer over-read. If exploited, this can lead to exposure of sensitive information – confidentiality loss.

//generated// Missing Code (in 'dtls1_process_heartbeat(SSL *s)') to Range Verify length (1 + 2 + 16 <= s->s3->rrec.length 1 + 2 + payload + 16 <= s->s3->rrec.length) Transferred (via network) in Third-Party (ssl /d1_both.c#L1462 ssl /t1_lib.c#L2591) Local leads to Inconsistent Value ('payload')

, which propagates to Wrong Size (in 'memcpy(bp, pl, payload)') Sequential Reposition (pointer) Heap Used (for s->s3->rrec.data[0]) Third-Party (ssl /d1_both.c#L1487 ssl /t1_lib.c#L2620) in Userland resulting in Over Bounds Pointer ('pl')

, which propagates to Over Bounds Pointer (in 'memcpy(bp, pl, payload)') Sequential Read (object) Huge (up to 64kb per exploit) Heap Used Third-Party (ssl /d1_both.c#L1487 ssl /t1_lib.c#L2620) in Userland resulting in Buffer Over-Read ('pl')

. If exploited this can lead to IEX (confidentiality loss).



Show/Hide Definitions

vendor:product: openssl:openssl

[Bug Report](#)

[Code with Bug](#)

[Code with Fix](#)

[NVD Entry](#)

Class

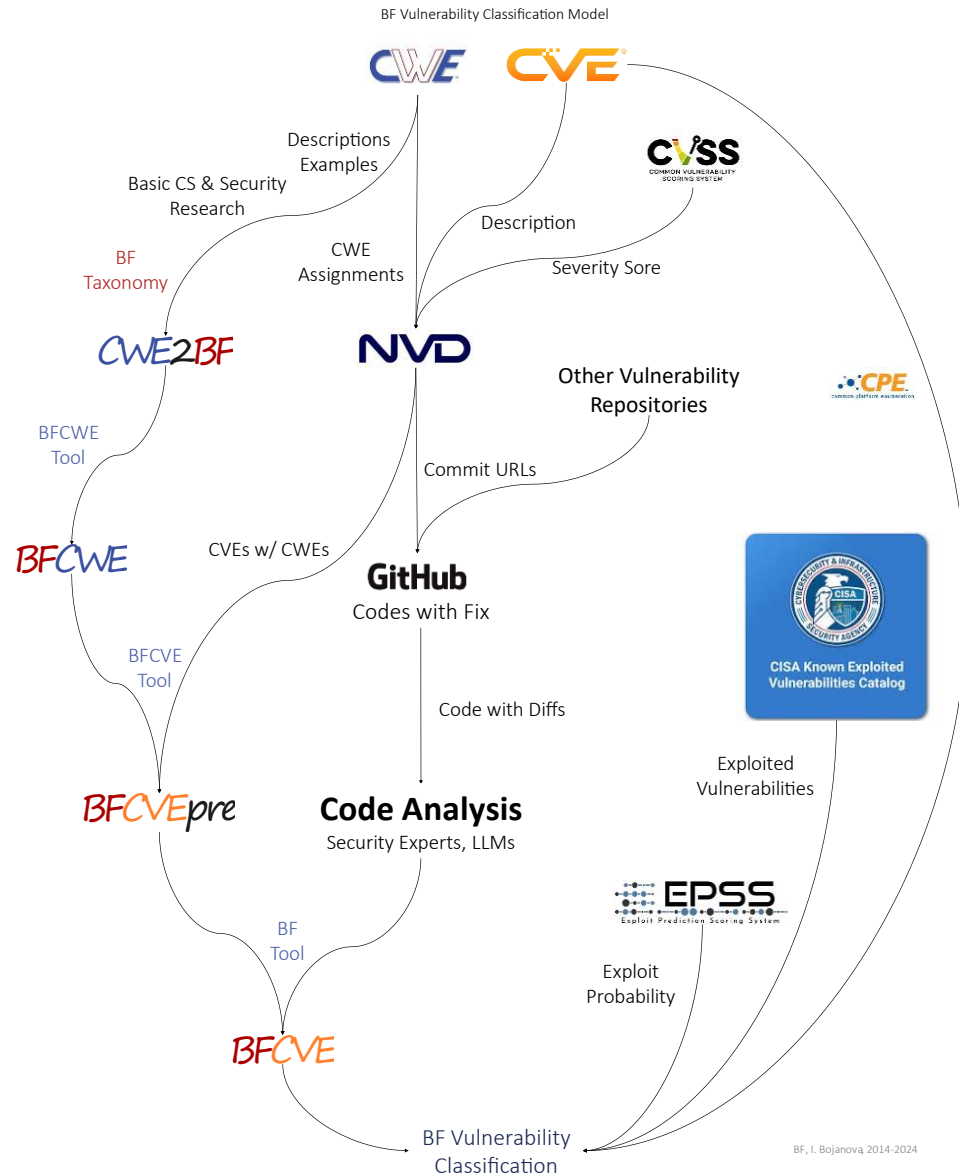
Definition

DVR

Data Verification (DVR) class – Data are verified (semantics check) or corrected (assign, remove) improperly.

...

BF Vulnerability Classification



```

]with cweClass as (
  select distinct c.Type, class = c.Name, wo.cwe
  from bf.class c
  inner join bf.operation o on c.Name = o.Class
  inner join cwebf.operation wo on o.Name = wo.operation
)
select m.cve [CVE], m.cwe [CWE], n.score [CVSS], ci.url [CodeWithFix], c.Type [BFClassType],
       c.class [BFClass], v.cause [Cause], v.operation [Operation], v.consequence [Consequence]
from cweClass c
inner join nvd.mapCveCwe m on m.cwe = c.cwe
inner join nvd.cve n on m.cve = n.cve
inner join gitHubVul.cve u on u.cve = n.cve
inner join gitHubVul.commitId ci on ci.id = u.commitId
inner join cwe.cwe w on w.id = m.cwe
inner join cwebf.specification s on s.cwe = m.cwe
inner join cwebf.mainWeakness mw on mw.mainWeakness = s.mainWeakness
inner join bf.validWeakness v on v.id = mw.weakness
left outer join cwebf.otherWeakness cw on cw.cwe = m.cwe and cw.mainWeakness = s.mainWeakness
left outer join bf.validWeakness vv on vv.id = cw.weakness
left outer join bf.operation oo on oo.Name = vv.operation
left outer join bf.class cc on oo.Class = cc.Name
where (c.Type = '_MEM')
order by n.score desc, m.cve, s.cwe, cw.chainId

```

NVD's One-to-Five Year Plan

Once the NVD is up and running, Brewer said the program will consider new approaches to improving its processes within the next one to five years, especially around software identification.

Some of the ideas include:

-
- ▶ **Involving more partners:** Being able to have outside parties submit CPE data for the CPE Dictionary in ways that scale to fit the ever-growing number of IT products

 - ▶ **Software identification improvements:** Dealing with software identification in the NVD in a way that scales with growing complexities (the adoption of PURLS is considered)

 - ▶ **New types of data:** Developing capabilities to publish additional kinds of data to the NVD (e.g. from EPSS, NIST Bugs Framework)

 - ▶ **New use cases:** Developing a way to make NVD data more consumable and more customizable to targeted use cases (e.g. getting email alerts from NVD when CVEs are published)

 - ▶ **CVE JSON 5.0:** Expanding the NVD's capabilities to utilize new data points available in CVE JSON 5.0

 - ▶ **Automation:** Developing a way to automate at least some CVE analysis activities
-

<https://www.infosecurity-magazine.com/news/nist-unveils-new-nvd-consortium/>

Machine readable formats of:

- BF taxonomy
- BFCWE specifications
- BFCVE specifications
- Vulnerability classifications

✓ Projects related to:

- Vulnerability specification generation
- Bug detection
- Vulnerability analysis and remediation
- Security failures and risks

Improving the Software Security **Triaging** and **Remediation** Processes using Hybrid Techniques along with Human-Readable Diagnoses

A dissertation by: Kedrian James
Committee: Fabian Monrose (advisor), Prasun Dewan, Cynthia Sturton, Sridhar Duggirala and Michalis Polychronakis

Approach

Report Generation Process

- Uses NIST's Framework¹
 - Taxonomic model for describing bugs.
 - Each class is comprised of a set of:
 - Causes, Operations, Consequences, Operands, Attributes, Sites.

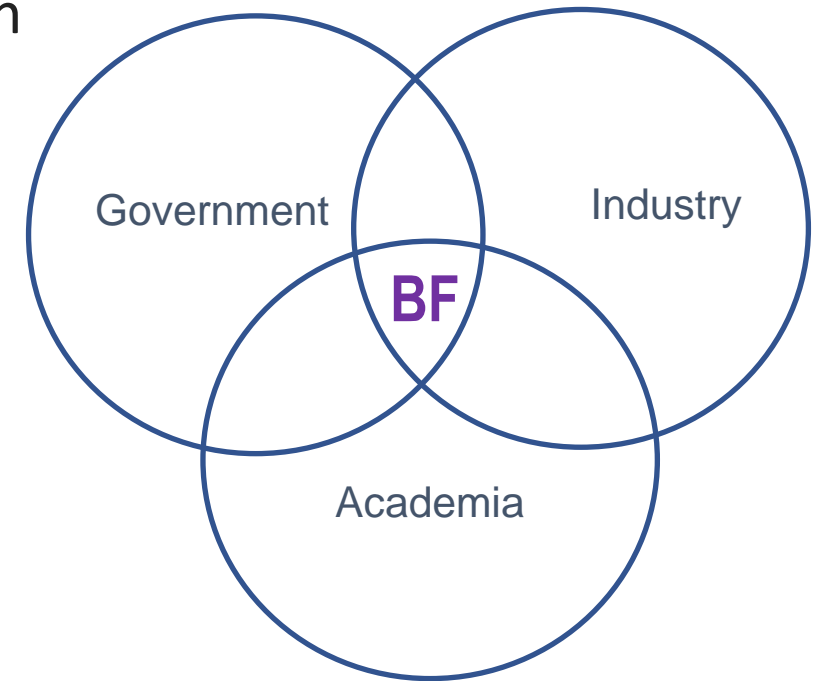
Challenge: Model is theoretical and lacks detail on how to collect information for crash diagnosis.

1. I. Bojanova, NIST, The Bugs Framework (BF), Accessed: 12/16/2023, [Online]. Available: <https://usnistgov.github.io/BF/>.

BF – Potential Impact

BF – Potential Impacts

- Allow precise communication about security bugs, weaknesses, and vulnerabilities
- ML/AI bug finding, vulnerability analysis, and resolution
- Help identify exploit mitigation techniques.



Questions

BF Contact



Irena Bojanova, BF PI & Lead

irena.bojanova@nist.gov



<https://samate.nist.gov/BF/>
<https://usnistgov.github.io/BF/>